Hybrid CP-SAT Solver for Temporal Planning and Scheduling

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Aries

New experimental solver¹

- for experimenting with hybrid SAT/CP solving
- targeting Temporal Planning

First milestone: a CP-SAT solver for Disjunctive Scheduling

- Core subproblem of temporal planning
- Not a strong suit of existing CP-SAT solvers

¹MIT licensed, in Rust: https://github.com/plaans/aries

Case Study: Disjunctive Scheduling

Given a set of **tasks** subject to **precedence** and **no-overlap** constraints, find a schedule of **minimum duration**.

Most famous: JobShop Scheduling Problem

JobShop: CSP model

Constants:

• $d_i \in \mathbb{N}$: duration of the i^{th} task.

(Decision) Variables:

• $s_i \in \mathbb{N}$: start time of the i^{th} task.

Constraints:

- precedence: $s_i + d_i \le s_j$
- no-overlap: $s_i + d_i \le s_j \ \lor \ s_j + d_j \le s_i \ \ \forall i, j \text{ s.t. } machine(i) = machine(j)$

Objective (makespan)

• Minimize $\max_i s_i + d_i$

 $\forall i, j \text{ s.t. } job(i) = job(j) \land i < j$

Aries: CP Solver Fundamentals

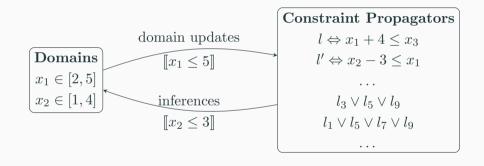
Domains

 $x_1 \in [2, 5]$

 $x_2 \in [1, 4]$

- Variables: (bounded) integers
- \bullet Domains: Lower & Upper Bounds
- Search events: Bound changes

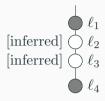
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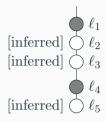


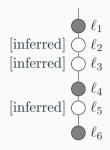


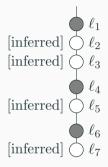


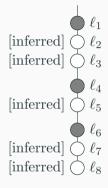


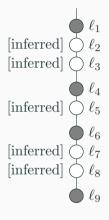


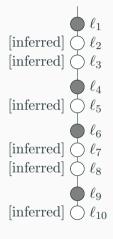


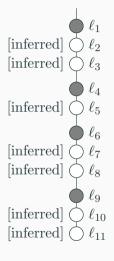


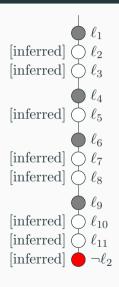




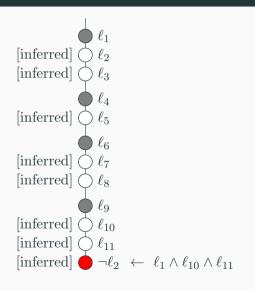




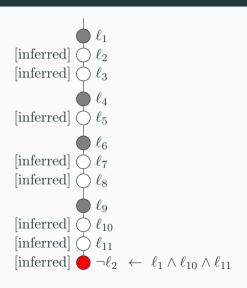




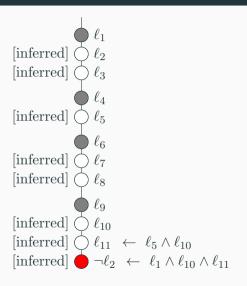
$$\ell_2 \wedge \neg \ell_2 \to \bot$$



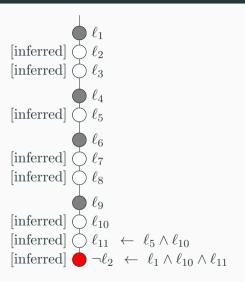
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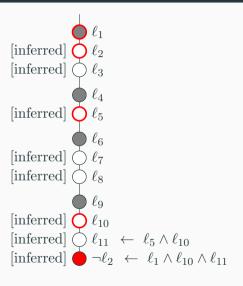
$$\begin{array}{c} \ell_2 \wedge \neg \ell_2 \to \bot \\ \\ \ell_2 \wedge (\ell_1 \wedge \ell_{10} \wedge \ell_{11}) \to \bot \end{array} \qquad \text{(resolved } \neg \ell_2\text{)} \end{array}$$



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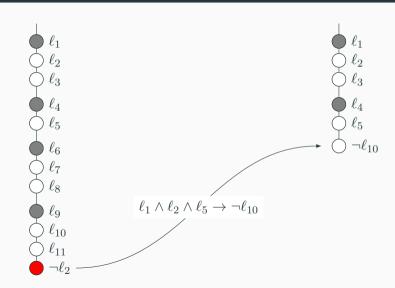
Conflict:

$$\begin{array}{c} \ell_2 \wedge \neg \ell_2 \to \bot \\ \ell_2 \wedge (\ell_1 \wedge \ell_{10} \wedge \ell_{11}) \to \bot & \text{(resolved } \neg \ell_2) \\ \ell_2 \wedge (\ell_1 \wedge \ell_{10} \wedge (\ell_5 \wedge \ell_{10}) \to \bot & \text{(resolved } \ell_{11}) \\ \ell_1 \wedge \ell_2 \wedge \ell_5 \wedge \ell_{10} \to \bot & \text{(reorganized)} \end{array}$$

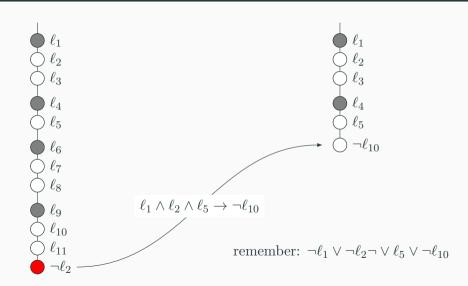
 ℓ_{10} : Unique Implication Point (UIP)
Asserting clause:

$$\ell_1 \wedge \ell_2 \wedge \ell_5 \to \neg \ell_{10}$$

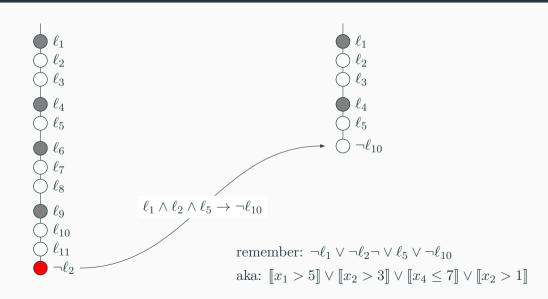
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- (lazily) create boolean variable for every bound literal
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To handle:
$$[x_1 > 5] \lor [x_2 > 3] \lor [x_4 \le 7] \lor [x_2 > 1]$$

- Create boolean variables
 - $\ell_1 : [x_1 > 5]$
 - $\ell_2 : [x_2 > 3]$
 - $\ell_3 : [x_4 \le 7]$
 - $\ell_4 : [x_2 > 1]$

- Post the disjunctive constraint
 - $\ell_1 \vee \ell_2 \vee \ell_3 \vee \ell_4$
- Maintain consistency of $\ell_1/x_1, \ell_2/x_2, \ell_3/x_4, \ell_4/x_2$

Aries: Reasons natively on bound literals

$$[x_1 > 5] \lor [x_2 > 3] \lor [x_4 \le 7] \lor [x_2 > 1]$$

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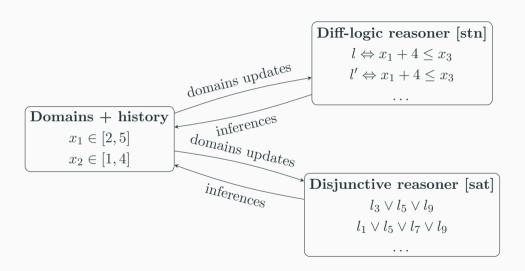
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- No variable creation
- No synchronization costs

Downside: pervasive change

Structure of Propagators: SMT-like "reasoners"



Difference-Logic Reasoner (aka DTN propagator)²

²SMT-like, organized propagation + few tricks

Search Strategy: Decision Variables

Which decision variables to branch on?

- $t_i \in \mathbb{N}$: start times
- $o_{ij} \in \{0,1\}$: task ordering
 - $[o_{ij} \ge 1] \Leftrightarrow t_i + d_i \le t_j$
 - $[o_{ij} \le 0] \Leftrightarrow t_j + d_j \le t_i$

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Key idea: maximize the number of conflicts per decision

 \hookrightarrow prefer literals that participate in conflicts

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- clause literals: $\ell_1, \ell_2, \ell_5, \ell_{10}$
- resolved literals: ℓ_2, ℓ_{11} (seen while building the clause)

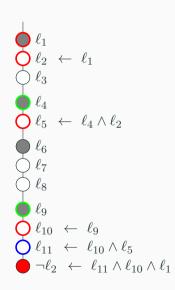


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- clause literals: $\ell_1, \ell_2, \ell_5, \ell_{10}$
- resolved literals: ℓ_2, ℓ_{11} (seen while building the clause)
- reasoned literals: ℓ_1, ℓ_4, ℓ_9 (reason of literals in the clause)



Search Strategy: full picture

- Variable selection: Learning Rate Branching (LRB)
 - Favors conflictual variables
- Value selection: solution guided
 - take value from last solution
 - search focused in its neighborhood
- Greedy Search initialization
 - until first conflict, assign start-time variables
 - Choice: Earliest-starting time, min-slack

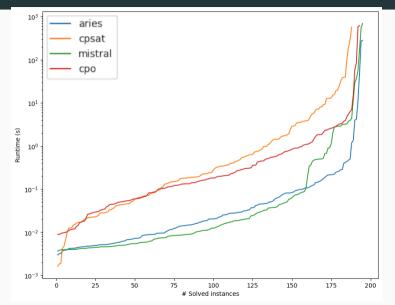
Performance Evaluation: Contenders

State-of-the-art CP solvers for disjunctive scheduling.

- Mistral: Pure CP for disjunctive scheduling
- CPOPTIMIZER: SOTA CP Scheduling solver
- CPSAT: SOTA Hybrid CP-SAT solver (OrTools)

Evaluated on common ${\bf OpenShop}$ and ${\bf JobShop}$ instances

OpenShop: runtimes (cactus plot)

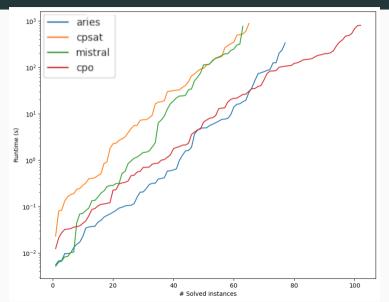


Runtime necessary to find proven optimal solutions.

ARIES & MISTRAL:

- All problems solved
- Aries systematically faster when more than 1 second is needed

Jobshop: runtimes (cactus plot)



Runtime necessary to find proven optimal solutions.

ARIES & CPOPTIMIZER:

- Aries fails to prove optimality for the harder instances
- Solution Quality difference is of 0.2%
- Aries generally substantially faster

Conclusion

Aries on disjunctive scheduling **milestone**:

- simple and generic design
- state-of-the-art performance
- extensible CP core

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... on the road to temporal planning

- resource reasoning
- optional reasoning

Aries for Temporal Planning

Structures in temporal planning problems

Action Template:

```
\begin{aligned} & \mathsf{move}(r, a, b) \\ & \mathsf{variables:} \quad r, \ a, \ b, \ t_{start}, \ t_{end} \\ & \mathsf{constraints:} \quad a \neq b \\ & \quad t_{end} - t_{start} = \mathit{travel-time}(a, b) \\ & \mathsf{conditions:} \quad [t_{start}] \ \mathit{loc}(r) = a \\ & \quad \mathsf{effects:} \quad ]t_{start}, t_{end}] \ \mathit{loc}(r) \leftarrow b \end{aligned}
```

From actions to CSP

For each action template:

- generate a bounded set of optional activities
- each with its own parameters (decision variables)

$$move^{1}(r^{1}, a^{1}, b^{1}) : p^{1}$$

 $move^{2}(r^{2}, a^{2}, b^{2}) : p^{2}$

• constraints ensure coherence of activities + goal achievement

Problem as a dummy action

```
Problem (C_0)
variables: t, \ell
constraints: t < 100
\ell = London \lor \ell = Dublin
conditions: [t, t] loc(Bob) = Toulouse \leftarrow end condition
effects: [0, 0] loc(Bob) \leftarrow Toulouse \leftarrow init. state
```

Decomposition validity: effect token

Effect token

$$move^{1}(r^{1}, a^{1}, b^{1})$$

$$\cdots$$
effects: $]t^{1}_{start}, t^{1}_{end}]loc(r^{1}) \leftarrow b^{1}$

$$\cdots$$

$$b^{1}_{start}$$

$$t^{1}_{end}$$

$$t^{1}_{end}$$

Here: t^1 is a lower bound on the persistence of the effect.

Decomposition validity: condition token

Condition token

$$move^{1}(r^{1}, a^{1}, b^{1}) : p^{1}$$

$$\dots$$
 conditions:
$$[t^{1}_{start}]loc(r^{1}) = a^{1}$$

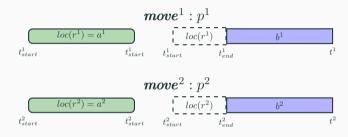
$$\dots$$

$$loc(r^{1}) = a^{1}$$

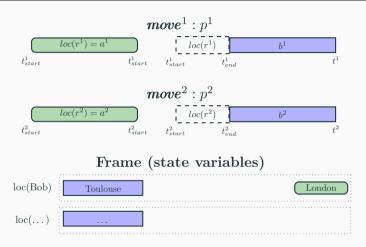
$$t^{1}_{start}$$

$$t^{1}_{start}$$

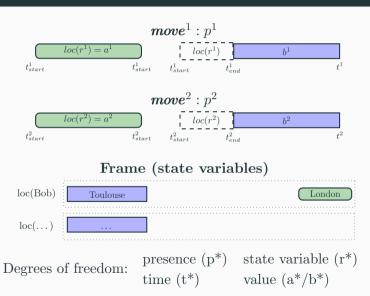
Planning Puzzle



Planning Puzzle

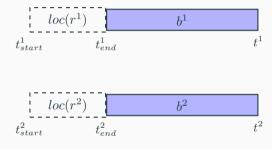


Planning Puzzle



Rule 1: No overlapping effects (coherence)

No two overlapping effects



$$\begin{array}{ccc} p^1 \wedge p^2 & \Longrightarrow & t^1 \leq t_{start}^2 \; \mathsf{V} \\ & & t^2 \leq t_{start}^1 \; \mathsf{V} \\ & & r^1 \neq r^2 \end{array}$$

Rule 2: All conditions are *supported* (support)

For any condition C^1 , there is an effect E^2 such that:

$$p^{2}$$

$$t_{end}^{2} \leq t_{start}^{1} \wedge t_{end}^{1} \leq t^{2}$$

$$r^{1} = r^{2}$$

$$a^{1} = b^{2}$$

$$loc(r^{1})$$
("each green must be in a blue")

Bounded planning problem as CSP

CSP (X, C) where

- X = all variables in chronicles + chronicle presence
 - $\{r^1, \ldots, p^1\} \cup \{r^2, \ldots, p^2\} \cup \{t, \ell\}$
- C = all constraints:
 - coherence
 - \bullet support
 - ...

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- C = all constraints:
 - coherence
 - support
 - ...
 - refinement (HTN) / symmetry breaking (generative)
 - resources (numeric)

 \Rightarrow handed over to Aries' CP/SAT solver

Aries planner

Planner features

- Generative or Hierarchical planning
- Durative actions
- Timed effects (TILs), timed goals (deadline)
- (Simple) Numeric Planning
- Optimization (length, costs, makespan, final value)

Integration in Unified Planning

- python library
- modeling/solving
- from AIPlan4EU project



Aries: at the state-of-the art?

- When action space is constrained (hierarchical non-recursive)
- temporally expressive (deadlines, time-windows)
- when quality matters

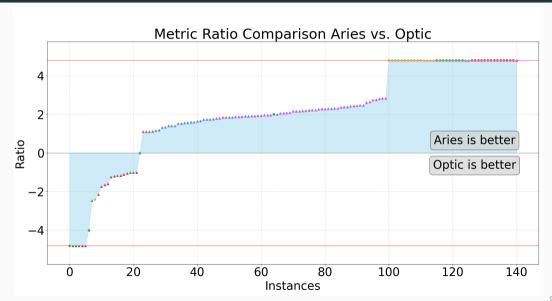
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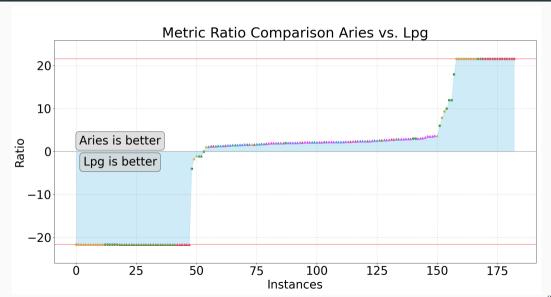
Hierarchical: no competitors Generative, compared to:

- Optic, (much) better coverage, higher quality
- LPG, lower coverage, higher quality

Aries vs Optic: solution quality (temporal+numeric problems)



Aries vs LPG: solution quality (temporal+numeric problems)



References

- Arthur Bit-Monnot. Enhancing Hybrid CP-SAT Search for Disjunctive Scheduling. ECAI 2023.
- Arthur Bit-Monnot. Experimenting with Lifted Plan-Space Planning as Scheduling: Aries in the 2023 IPC. 2023 International Planning Competition (2023)
- Roland Godet, Arthur Bit-Monnot. Chronicles for Representing Hierarchical Planning Problems with Time. ICAPS Hierarchical Planning Workshop (HPlan 2022),